On the strong persistence property for monomial ideals
by
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Abstract

Let \( I \) be the edge ideal associated to a graph with loops, a weighted graph or a clutter. In this paper we study when \( I \) has the strong persistence property, this is \((I^{k+1}:I) = I^k\) for each \( k \geq 1 \).

Key Words: Strong persistence property, monomial ideals, clutters, weighted graphs, graphs with loops.

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1 Introduction

Let \( R \) be a commutative Noetherian ring. The associated primes set of an ideal \( I \) is \( \text{Ass}(I) = \{ P \in \text{Spec}(R) \mid P = (I:a) \text{ for some } a \in R \} \). If \( I = Q_1 \cap \cdots \cap Q_s \) is a minimal primary decomposition of \( I \), then \( \text{Ass}(I) = \{ r(Q_1), \ldots, r(Q_s) \} \) where \( r(Q_i) \) is the radical of \( Q_i \). \( I \) has the persistence property if \( \text{Ass}(I^k) \subseteq \text{Ass}(I^{k+1}) \) for each \( k \). In [2] it is showed that the edge ideal of a simple graph has the persistence property, and they use that these edge ideals satisfy \((I^{k+1}:I) = I^k\) for each \( k \). Recently was proved that this concept implies the persistence property (see [1]) and it is called the strong persistence property. These concepts are not equivalent, in [2, Example 2.18] is given a squarefree monomial ideal with the persistence property, but it does not have the strong persistence property. Assuming this terminology, in [2, Lemma 2.12] was proved that the edge ideal of a simple graph has the strong persistence property. In this paper we study the strong persistence property for edge ideals of graphs with loops, weighted graphs, and clutters.

This paper is organized as follows: in Sect. 2 we prove that the edges ideals of graphs with loops have the strong persistence property. In Sect. 3 we prove that the edge ideal of a vertex–weighted graph \((G, w)\) has the strong persistence property. Furthermore, we prove that \( I(G)^k \) and \( I(G, w)^k \) have the same associated primes. In Sect. 4 we study the edges ideals of clutters. In particular, we show that a clutter has the strong persistence property if and only if at least one of its connected components has the strong persistence property. Also, we prove that a König unmixed clutter without 4-cycles and squarefree monomials in four variables have the strong persistence property. Furthermore, we show that \((I^2: I) = I\), if \( I \) is a squarefree monomial ideal. Finally we prove that the strong persistence property is closed under \( c \)-minor and cones. In Sect. 5 we give some properties of the strong persistence property. Also, we introduce the symbolic persistence property and we show that an ideal has this property if it has the strong persistence property.
2 Graphs with loops

A graph with loops is a triplet $G = (V, E, L)$ where $G = (V, E)$ is a simple graph with $V = \{x_1, \ldots, x_n\}$ and $E \subseteq \{(x_i, x_j) \mid x_i \in V\}$. $L$ is called the set of loops of $G$. Let $R = K[x_1, \ldots, x_n]$ be a polynomial ring, as usual we use $x^a$ as abbreviation for $x_1^{a_1} \cdots x_n^{a_n}$, where $a = (a_1, \ldots, a_n)$ is an integer vector with $a_i \geq 0$. If $f = (x_i, x_j) \in E$ or $f = (x_i, x_i) \in L$, then we take the monomial $f = x_i x_j$ or $\tilde{f} = x_i^2$, respectively. The edge ideal of a graph with loops $G = (V, E, L)$ is the ideal $I(G) = \langle \{f_i \mid f_i \in E \cup L\} \rangle = I(G) + \langle \{x_i^2 \mid (x_i, x_i) \in L\} \rangle$ where $I(G) = \langle \{(x_i, x_j) \mid (x_i, x_j) \in E\} \rangle$ is the edge ideal of $G = (V, E)$.

**Example 1.** Graph with loops, where $L = \{(x_1, x_1), (x_3, x_3)\}$.

```
    x1 --- x2
     |    |    \
     v    v    v
    x3 --- x4
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For an integer vector $a = (a_1, \ldots, a_n)$ with $a_i \geq 0$, we define the simple graph $G^a$ with vertex set is $V^a = \{x_1^1, \ldots, x_n^1, x_1^2, \ldots, x_n^2, \ldots, x_1^n, \ldots, x_n^n\}$, and whose edge set

$$E^a = \left\{ \{x_i^{k_i}, x_j^{k_j}\} \mid \begin{array}{l}
(x_i, x_j) \in E, k_i \leq a_i, \text{ and } k_j \leq a_j; \text{ or } \\
(x_i, x_j) \in L \text{ and } 1 \leq k_i < k_j \leq a_i
\end{array} \right\},$$

where $j$ in $x_i^j$ is only an index. Furthermore, if $x_i \in V(G)$, then we define the duplication of $x_i$ in $G^a$ as the simple graph $G^{a_{x_i}} = G^{a+e_i}$ where $e_i$ is the $i$-th unit vector in $\mathbb{R}^n$. This operation is commutative, that is $(G^{a_{x_i}})^{x_j} = (G^{a_{x_j}})^{x_i}$ for each $x_i, x_j \in V$. Furthermore, if $f = (x_i, x_j) \in E$, then we denote by $(G^a)^f = G^{a+e_i+x_j}$; and if $f = (x_i, x_i) \in L$, then $(G^a)^f = G^{a+e_i+x_i}$. 

**Definition 1.** Let $G$ a simple graph. A matching of $G$ is a set of pairwise disjoint edges. The matching number of $G$, denoted by $\nu(G)$, is the size of any maximum matching of $G$. A matching that covers all vertices of $V(G)$ is called a perfect matching of $G$.

**Notation 1.** $\text{Mon}(R)$ is the set of monomials in $R = K[x_1, \ldots, x_n]$. If $I = (m_1, \ldots, m_s)$ with $m_i \in \text{Mon}(R)$, then $G(I)$ is the minimal monomial generating set of $I$.

**Proposition 1.** Let $G$ be a graph with loops whose vertex set is $V = \{x_1, \ldots, x_n\}$. If $a = (a_1, \ldots, a_n)$ is an integer vector where $a_i \geq 0$, then $G^a$ has a matching of size $l$ if and only if $x^a \in I(G)^l$.

**Proof.** (⇒) Let $P = \{g_1, \ldots, g_l\}$ be a matching of $G^a$ where $g_j = \{x_i^{k_i}, x_j^{s_j}\}$. Now, we consider the monomial $x^b = \prod_{i=1}^{l} x_i^{k_i} x_j^{s_j} \in I(G)^l$. If $b = (b_1, \ldots, b_n)$, then $b_i = |\{r \mid x_r^{r_i} \in \bigcup_{j=1}^{l} g_j\}|$ for each $1 \leq i \leq n$. Since $P$ is a matching, $b_i \leq a_i$. Therefore, $x^b | x^a$ and $x^a \in I(G)^l$.

(⇐) We take $E \cup L = \{f_1, \ldots, f_q\}$. If $x^a \in I(G)^l$, then there exist an integer vector $\alpha = (\alpha_1, \ldots, \alpha_q)$ such that $\alpha_1 + \cdots + \alpha_q = l$ and $x^a = m \tilde{f}_1^{a_1} \cdots \tilde{f}_q^{a_q}$ with $m \in \text{Mon}(R)$. We can assume that $\alpha_1 > 0$. If $f_1 = \{x_r, x_s\} \in E$, then $\alpha_1 = a_r$ and $\alpha_1 = a_s$ since $\tilde{f}_1^{a_1} | x^a$. If
$P_1 = \{g_1, \ldots, g_{\alpha_1}\}$, where $g_j = \{x_{r^2-j+1,} x_{r^2-j+1}^2\}$ for $j \leq \alpha_1$, then $P_1$ is a matching of $G^a$ of size $\alpha_1$. If $f_1 = (x_r, x_r) \in L$, then $2\alpha_1 \leq a_r$ since $f_1^2 | x^a$. Consequently, $P_1 = \{g_1, \ldots, g_{\alpha_1}\}$, where $g_j = \{x_{r^2-j+2,} x_{r^2-j+2}^2\}$ for $j \leq \alpha_1$, is a matching of $G^a$ of size $\alpha_1$. Hence,

$$G^b = G^a \cup g_j$$

and

$$x^b = \frac{x^a}{f_1^2} = m f_2^2 \cdots f_{\alpha_q}^q \in I(G)^{l-\alpha_1},$$

where $b = a - \alpha_1(e_r + e_s)$ if $f_1 \in E$ or $b = a - 2\alpha_1 e_r$ if $f_1 \in L$. Following with the processes, we obtain matchings $P_1, \ldots, P_q$ such that

$$V(P_{i+1}) \cap \bigcup_{j=1}^{i} V(P_j) = \emptyset$$

since $V(P_{i+1}) \subseteq V(G^a) \setminus \bigcup_{j=1}^{i} V(P_j).$

Therefore, $\bigcup_{j=1}^{q} P_j$ is a matching of $G^a$ of size $l$. \qed

**Corollary 1.** $x^a \in I(G)^k \setminus I(G)^{k+1}$ if and only if $k = \nu(G^a)$.

**Definition 2.** The deficiency of a simple graph $G$ is given by

$$\text{def}(G) = |V(G)| - 2\nu(G),$$

where $c_0(G)$ denotes the number of odd components (components with an odd number of vertices) of $G$.

**Theorem 1.** ([2]) If $G$ is a simple graph, then

$$\text{def}(G) = \max\{c_0(G \setminus S) - |S| \mid S \subseteq V(G)\},$$

where $c_0(G)$ denotes the number of odd components of $G$. We obtain that the odd connected components of $G^a$ for $k = 1, \ldots, h_r$ are $H_1, H_2, \ldots, H_{r-1}, H_{r+1}, \ldots, H_r$. Consequently,

$$c_0(G^a) - |S| > \delta = \text{def}(G^a).$$

A contradiction. Now, assume $f = \{x_i, x_j\} \in E(G)$. If $x_i, x_j \in E(H_k)$, then we consider the subgraph $H'_k$ of $G^a \setminus S$ induced by $V(H_k) \cup \{x_i^{a_{r+1}}, x_j^{a_{r+1}}\}$. We obtain that the odd connected components of $G^a$ for $k = 1, \ldots, h_r$ are $H_1, H_2, \ldots, H_{r-1}, H_{r+1}, \ldots, H_r$. So,

$$c_0(G^a) - |S| > \delta = \text{def}(G^a).$$
This implies $V(H_k) = \{x_i^k\}$ and $a_j = 0$ or $x_j^{k_j} \in S$ for each $k_j \leq a_j$. Hence, the odd components of $G^{af} \setminus (S \cup \{x_j^{a_j+1}\})$ are $H_1, \ldots, H_r, \{x_i^{a_i+1}\}$. Thus,

$$c_0(G^{af} \setminus (S \cup \{x_j^{a_j+1}\})) - |S \cup \{x_j^{a_j+1}\}| = c_0(G^a \setminus S) - |S| > \delta = \text{def}(G^{af}).$$

A contradiction, therefore $\text{def}(G^a) = \text{def}(G^{af})$ for all $f \in F$. Therefore, $\nu(G^{af}) = \nu(G^a) + 1$, since $|V(G^{af})| = |V(G^a)| + 2$ for all $f \in F$.

\[ \begin{align*}
\Rightarrow & \; \text{def}(G^{af}) = |V(G^{af})| + 2\nu(G^{af}) = |V(G^a)| + 2 - 2(\nu(G^a) + 1) = \text{def}(G^a) = \delta. \\
\end{align*} \]

\[ \begin{proof}
\end{proof} \]

**Theorem 2.** $I(G)$ has the strong persistence property if $G$ is a simple graph.

\[ \begin{proof}
\end{proof} \]

**Theorem 3.** If $G$ is a graph with loops, then $(I^{k+1}:I) = I^k$ with $I = I(G)$.

\[ \begin{proof}
\end{proof} \]

**Corollary 2.** $I(G)$ has the persistence property if $G$ is a graph with loops.

\[ \begin{proof}
\end{proof} \]

3 Weighted monomial ideals

Let $I$ be a monomial ideal, an irreducible monomial ideal $J$ is $I$-minimal if $J$ is minimal in the set of irreducible monomial ideals (have the form $\{x_i^{\alpha_i}, \ldots, x_s^{\alpha_s}\}$), such that $I \subseteq J$. The set of $I$-minimal ideals is a minimal primary decomposition of $I$.

**Definition 3.** For $m_1, m_2 \in \text{Mon}(R)$, $m_1 \parallel m_2$ if $m_1 | m_2$ and $m_1^{s+1} \not\parallel m_2$.

**Proposition 3.** Let $I$ be a monomial ideal. If $\{x_1^{\alpha_1}, \ldots, x_s^{\alpha_s}\}$ is an $I$-minimal ideal, then for each $1 \leq t \leq s$ there is $m \in G(I)$ such that $x_i^{\alpha_i} | m$.

\[ \begin{proof}
\end{proof} \]

Since $J = \{x_1^{\alpha_1}, \ldots, x_s^{\alpha_s}\}$ is an $I$-minimal ideal, then $I \subseteq J$. Thus, if $x_i^{\alpha_i} \not\parallel u$ for each $u \in G(I)$, then $I \subseteq (\{x_1^{\alpha_1}, \ldots, x_s^{\alpha_s}\} \setminus \{x_i^{\alpha_i}\})$. This contradicts the minimality of $J$. Hence, $x_i^{\alpha_i} | u$ for some $u \in G(I)$. Now, if $x_i^{\alpha_i+1} | m$ for each $m \in G(I)$ such that $x_i^{\alpha_i} | m$, then $I \subseteq (\{x_1^{\alpha_1}, \ldots, x_i^{\alpha_i+1}, \ldots, x_s^{\alpha_s}\}) \not\subseteq J$. A contradiction, therefore there is $m \in G(I)$ such that $x_i^{\alpha_i} | m$.\[ \begin{proof}
\end{proof} \]
Definition 4. A weight over a polynomial ring $R = K[x_1, \ldots, x_n]$ is a function $w : \{x_1, \ldots, x_n\} \to \mathbb{N}$, $w_i = w(x_i)$ is called the weight of the variable $x_i$. Given a monomial ideal $I$ and a weight $w$, the weighted ideal of $I$ and $w$ is $I_w = \{h(m) \mid m \in G(I)\}$ where $h$ is the isomorphism $h : R \to K[x_1^{w_1}, \ldots, x_n^{w_n}]$ given by $x_i \mapsto x_i^{w_i}$.

Remark 1. Since $h$ is an isomorphism, $G((I_w)^k) = G((I^k)_w)$, so $(I_w)^k = (I^k)_w$.

Theorem 4. Let $I$ be a monomial ideal and $w$ a weight over $R$, then

i) $\text{Ass}(I_w^k) = \text{Ass}(I^k)$ for each $k$;

ii) $I$ has the persistence property if and only if $I_w$ has the persistence property;

iii) $I$ has the strong persistence property if and only if $I_w$ has the strong persistence property.

Proof. i) If $(x_1^{\beta_1}, \ldots, x_n^{\beta_n})$ is an $I^k$-minimal ideal, then by Proposition 3 there is $m' \in G(I^k)$ such that $x_{i_j}^{\beta_i} \mid m'$, so there is $r_i$ such that $\beta_i = w_i r_i$ for $1 \leq j \leq s$. If $m = x_1^{\alpha_1} \cdots x_n^{\alpha_n} \in G(I^k)$, then $h(m) = x_1^{\alpha_1 w_i} \cdots x_n^{\alpha_n w_n} \in G(I^k) \subseteq (x_1^{\beta_1}, \ldots, x_n^{\beta_n})$. Hence, there exist $t \leq s$, such that $x_{i_t}^{\beta_t} \mid h(m)$. Thus, $h_r = h_i \leq w_i r_i$ implies $h_r \leq w_i r_i$ and $x_{i_t}^{\beta_t} \mid m$. Consequently $I^k \subseteq (x_1^{r_1}, \ldots, x_n^{r_n})$. Now, if $(x_1^{\alpha_1}, \ldots, x_n^{\alpha_n})$ is an $I^k$-minimal, then $I^k \subseteq (x_1^{\alpha_1}, \ldots, x_n^{\alpha_n})$. So, $(x_1^{r_1}, \ldots, x_n^{r_n})$ is $I^k$-minimal and $(x_1^{w_1}, \ldots, x_n^{w_n})$ is $I^k$-minimal. Therefore, $(x_1^{\alpha_1}, \ldots, x_n^{\alpha_n})$ is $I^k$-minimal. Taking radicals of the $I^k$-minimal and $I^k$-minimal ideals we obtain $\text{Ass}(I^k) = \text{Ass}(I^k)$.

ii) By i).

iii) $\Rightarrow$) Since $h$ is an isomorphism of $k$-algebras between $R' = K[x_1^{w_1}, \ldots, x_n^{w_n}]$ and $R$, $h(I)$ has the strong persistence property in $K[x_1^{w_1}, \ldots, x_n^{w_n}]$. Also, $m = x_1^{a_1} \cdots x_n^{a_n} \in R'$ if and only if $w_1 \mid \lambda_i$ for each $i$. Thus, $I_w \cap [x_1^{w_1}, \ldots, x_n^{w_n}] = h(I)$. Now, if $m \in (I^k)_{w_1} \cap I_w$, then $gm = f_1 \cdots f_{k+1}$ for each $g \in G(I_w)$ where $g_i \in G(I_w)$. We take $m = c_1 \cdots c_n$ and $\ell = c_1^{\beta_1} \cdots c_n^{\beta_n}$. If $r_i$ and $t_i$ are the remainders obtained by dividing $a_i$ and $b_i$ by $w_i$ respectively, then $w_i \mid r_i - t_i$, since $G(I_w) \subseteq [x_1^{w_1}, \ldots, x_n^{w_n})$. So, $r_i = t_i$ and we take $m' = x_1^{\beta_1} \cdots x_n^{\beta_n}$ and $\ell' = c_1^{\beta_1} \cdots c_n^{\beta_n}$. Hence, $m' \ell' \in [x_1^{w_1}, \ldots, x_n^{w_n}]$ and $gm' = f_1 \cdots f_{k+1}$. Since $G(I_w) = G(h(I))$, $m' \in (h(I))^{k+1} = h(I) = h(I^{k+1})$ implies $m' \in I^k_w$. Therefore $m \in I^k_w$, since $m' \mid m$.

$\Leftarrow$) We take $m \in (I^{k+1}) \cap \text{Mon}(R)$, then $m f = \ell f_1 \cdots f_{k+1}$ with $f, f_1, \ldots, f_{k+1} \in G(I)$. So $h(m)h(f) = \ell h(f_1) \cdots h(f_{k+1}) \in I^{k+1}$. Thus, $h(m) \in (I^{k+1}) \cap I^k_w$ since $G(I_w) = G(h(I))$. This implies, $h(m) \in I^k_w \cap I^{k+1}$. Since $h(m) \in R' = K[x_1^{w_1}, \ldots, x_n^{w_n}]$, $\ell \in R'$. Therefore, $m = h^{-1}(\ell)g_1 \cdots g_{k+1} \in I^k$, since $h$ is an isomorphism.

Definition 5. A weighted graph $(G, w)$ consists of a simple graph $G$ and a function $w : V(G) \to \mathbb{N}$. The weight of $x \in V(G)$ is $w(x)$.

Definition 6. The edge ideal of the weighted graph $(G, w)$ denoted by $I(G, w)$ is the ideal generated by $\{x^{w_1} y^{w_2} \mid x, y \in E(G)\}$, where $w_k = w(x_k)$. 
Corollary 3. If $I = I(G)$ and $J = I(G,w)$, then $\text{Ass}(J^k) = \text{Ass}(I^k)$ for all $k$.

Proof. By Theorem 4, since $J = I_w$. □

Theorem 5. The edge ideal $I(G,w)$ has the strong persistence property.

Proof. By Theorem 4, since $I(G)$ has the strong persistence property. □

4 Squarefree monomial ideal

Let $R = K[x_1,\ldots,x_n]$ be a polynomial ring. A monomial ideal $I$ is squarefree if $G(I)$ consists of squarefree monomials. A clutter is a pair $C = (V,E)$ where $V$ is a finite set and $E$ is a set of subsets of $V$ such that if $a \subseteq b$ with $a, b \in E$, then $a = b$. The sets $V = V(C)$ and $E = E(C)$ are called vertex set and edge set, respectively. If $f = \{x_{i_1},\ldots,x_{i_r}\} \in E(C)$, then we denote by $\tilde{f}$ the squarefree monomial $x_{i_1}\cdots x_{i_r}$. Hence, if $f_1 \subseteq f_2 \subseteq X = \{x_1,\ldots,x_n\}$, then $f_1 \mid f_2$. The edge ideal of the clutter $C$, denoted by $I(C)$, is the ideal generated by $\{\tilde{f} \mid f \in E(C)\}$. This assignment defines a natural bijection between squarefree monomial ideals of $K[x_1,\ldots,x_n]$ and clutters whose vertex set is $X$. Finally we say that a clutter $C$ has the strong persistence property if $I(C)$ has the strong persistence property.

Lemma 1. Let $f, g$ be squarefree monomials, if there exists an integer $k \geq 2$ such that $f^k \mid mg$, then $f^{k-1} \mid m$.

Proof. Since $f^k \mid mg$, $mg = f^k \ell$ with $\ell \in \text{Mon}(R)$. We take $m' = \gcd(f,g)$, then $f = m'f'$ and $g = m'g'$ with $\gcd(f',g') = 1$. Hence, $\gcd(f,g') = \gcd(m',g') = u$. Consequently $u^2 \mid g$. But $g$ is a squarefree monomial, so $\gcd(f,g') = 1$. Thus $g' \mid \ell$, since $mg' = f'f^{k-1}\ell$. Therefore $m = f^{k-1}(f'u')$ where $\ell = u'g'$ implies $f^{k-1} \mid m$. □

Corollary 4. Let $I$ be a squarefree monomial ideal. If $G(I)$ has at most two elements, then $I$ has the strong persistence property.

Proof. Let $m$ be a monomial in $(I^{k+1}:I)$. So, for each $f \in I(G)$ there are monomials $\ell, g_1,\ldots,g_{k+1}$ with $g_i \in G(I)$, such that $mf = \ell g_1\cdots g_{k+1}$. If $f = g_i$ for some $i$, then $m \in I^k$. Now, if $f \neq g_i$ for each $i$, then $g_i = g_1$ since $|G(I)| \leq 2$. Thus $g_1^{k+1} \mid mf$. Hence, by Lemma 1, $g_1^k \mid m$ and $m \in I^k$. □

Theorem 6. If $I$ is a squarefree monomial ideal, then $(I^2:I) = I$.

Proof. Let $m$ be a monomial in $(I^2:I)$, then for each $f_1 \in G(I)$ there are $h_1, g_1 \in G(I)$ and a monomial $\ell_1$ such that $mf_1 = \ell_1 g_1 h_1$. Consequently,

$$m^2 f_1 = \ell_1 (mg_1) h_1 = \ell_1 \ell_2 g_2 h_2 h_1$$

where $mg_1 = \ell_2 g_2 h_2$ and $g_2, h_2 \in G(I)$. Follows, multiplying by $m$ we obtain

$$m^r f_1 = \ell_1\cdots\ell_r g_r h_r \cdots h_2 h_1,$$
where \( mg_{i-1} = \ell_i g_i h_i \) and \( g_i, h_i \in G(I) \) for \( 2 \leq i \leq r \). If \( r \geq |G(I)| \), then \( g_r = h_j \) or \( h_j = h_i \) for some \( 1 \leq i < j \leq r \). Hence, \( h_j^2 \mid m^r f_1 \) and by Lemma 1, \( h_j \mid m^r \). Thus, \( h_j \mid m \), since \( h_j \) is squarefree. Therefore \( m \in I \).

**Corollary 5.** If \( I \) is a squarefree monomial ideal and \( k \geq 2 \), then \( (I^k : I) \subseteq I \).

**Proof.** By Theorem 6, \( (I^k : I) \subseteq I \). Hence, \( I^k \subseteq I^2 \) and \( (I^k : I) \subseteq (I^2 : I) \). \( \square \)

**Theorem 7.** A clutter has the strong persistence property if and only if some of its connected components has the strong persistence property.

**Proof.** Let \( C_1, \ldots, C_r \) the connected components of \( C \) with \( V_i = V(C_i) \).

\( \Leftarrow \) We can suppose that \( C_1 \) has the strong persistence property. We take a monomial \( m \in (I^{k+1} : I) \). We can write \( m = m_1 \cdots m_r \) where \( m_i \in \text{Mon}(K[V_i]) \) and we take \( a_i \) such that \( m_i \in I^{s_i}_i \setminus I^{s_i + 1}_i \). For each \( f \in C_1 \) we consider \( s_f \) such that \( m_1 f \in I^{s_f}_1 \setminus I^{s_f + 1}_1 \) and \( s_1 = \min\{s_f \mid f \in C_1\} \). Thus, \( m_1 f \in I^{s_1}_1 \) for each \( f \in C_1 \), so \( m_1 \in (I^{s_1}_1 : I_1) = I^{s_1 - 1}_1 \). Hence,

\[
\begin{align*}
m &\in I^{s_1 + \sum a_i}_1 \quad \text{and} \quad m f \in I^{s_f + \sum a_i}_1 \setminus I^{s_f + 1 + \sum a_i}_1 \quad \text{for each} \quad f \in C_1,
\end{align*}
\]

Since \( m f \in I^{k+1} \), \( s_f + \sum a_i \geq k + 1 \). Then, \( s_1 + \sum a_i \geq k + 1 \). Therefore \( m \in I^k \).

\( \Rightarrow \) If \( I_i = I(C_i) \) has no the strong persistence property, then there is \( k_i \) and a monomial \( m_i \in (I^{k_i+1}_i : I_i) \setminus I^{k_i}_i \). We take \( a_i \) such that \( m_i \in I^{k_i}_i \setminus I^{k_i + 1}_i \), then \( a_i \leq k_i - 1 \). Now, we consider \( m = m_1 \cdots m_r \), then \( m \in I^b \setminus I^{b+1} \), for \( b = \sum a_i \). If we take \( f_i \in E(C_i) \), then \( m f_i \in I^{s_i} \), where \( s_i = a_1 + \cdots + k_i + 1 + \cdots + a_r \). But \( s_i \geq \sum a_j + 2 \), thus \( s = \min\{s_1, \ldots, s_r\} \geq \sum a_j + 2 \). Therefore \( m \in (I^s : I) \setminus I^{s-1} \).

**Example 2.** Let \( C \) be a clutter. If \( f_1, f_2 \in \{ A \subseteq V(G) \mid A \cap f = \emptyset \text{ if } f \in E(C) \} \), then by Theorem 7 and Corollary 4, \( C \cup \{ f_1, f_2 \} \) has the strong persistence property.

**Lemma 2.** Let \( C \) be a clutter. If there exists an edge \( f \in E(C) \) such that \( A = \{ g \cap f \mid g \in E(C) \} \) is a chain, then \( I(C) \) has the strong persistence property.

**Proof.** If \( m \) is a monomial in \( (I^{k+1} : I) \), then \( m \tilde{f} = \tilde{g} f_1 \cdots \tilde{f}_{k+1} \) where \( f_i \in E(C) \) and \( g \subseteq V(C) \). So, \( f \subseteq g \cup f_1 \cup \cdots \cup f_{k+1} \). Since \( A \) is a chain, we can assume \( f_{k+1} \cap f \subseteq f_k \cap f \subseteq \cdots \subseteq f_1 \cap f \). Thus, \( f \subseteq g \cup f_1 \) and \( f \mid \tilde{g} f_1 \). Therefore \( m \in I^k \).

**Corollary 6.** If \( C \) is a clutter without the strong persistence property, then for \( f \in E(C) \) there are \( f_1, f_2 \in E(C) \) such that \( f \cap f_1 \not\subseteq f \cap f_2 \) and \( f \cap f_2 \not\subseteq f \cap f_1 \).

**Definition 7.** Let \( C \) be a clutter. \( A \subseteq V(C) \) is a vertex cover if \( A \cap e \neq \emptyset \) for each \( e \in E(C) \). The cover number of \( C \) is \( \tau(C) = \min\{|A| \mid A \text{ is a vertex cover} \} \). \( C \) is unmixed if \( |B| = \tau(C) \) for each minimal vertex cover \( B \). A matching is a set of disjoint edges \( \{e_1, \ldots, e_s\} \) of \( C \). It is perfect if \( \bigcup_{i=1}^s e_i = V(C) \). Furthermore, \( C \) is König if there is a matching with \( \tau(C) \) edges.
Proposition 4. Let \( \mathcal{C} \) be a König clutter, then \( \mathcal{C} \) is unmixed if and only if there is a perfect matching \( e_1, \ldots, e_s \) with \( g = \tau(\mathcal{C}) \), such that for any two edges \( e \neq e' \) and for any two distinct vertices \( x \in e, y \in e' \) contained in some \( e_i \), one has that \( (e \setminus \{x\}) \cup (e' \setminus \{y\}) \) contains an edge.

Proof. See Corollary 2.11 in [3].

Definition 8. The incidence matrix of a clutter \( \mathcal{C} \), denoted by \( A_{\mathcal{C}} \), is the matrix whose columns are the characteristic vectors of the edges of \( \mathcal{C} \). A \( r \)-cycle of \( \mathcal{C} \) is a \( r \times r \)-submatrix of \( A_{\mathcal{C}} \) with exactly two 1’s in each row and each column.

Theorem 8. Let \( \mathcal{C} \) be a König unmixed clutter. If \( \mathcal{C} \) does not contain 4-cycles, then \( \mathcal{C} \) has the strong persistence property.

Proof. By Proposition 4, \( \mathcal{C} \) has a perfect matching \( e_1, \ldots, e_s \) where \( s = \tau(\mathcal{C}) \). If \( \mathcal{C} \) does not have the strong persistence property, then by Corollary 6 there exist \( f \in E(\mathcal{C}) \) and vertices \( x_1, x_2 \in (f, e_1 \setminus f_2) \) and \( x_2 \in (f_2, e_1) \setminus f_1 \). Now by Proposition 4, there exist \( f \in E(\mathcal{C}) \) such that \( f \subseteq (f \setminus e_1) \cup (f_2 \setminus x_2) \). We can assume \( e_1 \cap (f_2 \cup f_1) \) is minimal in

\[
B = \{ e_1 \cap (g_2 \cup g_1) \mid g_1, g_2 \in E(\mathcal{C}), g_2 \cap e_1 \notin g_1 \cap e_1, \text{ and } g_1 \cap e_1 \notin g_2 \cap e_1 \}.
\]

Thus, \( (e_1, f) \subseteq e_1 \cap ((f_1 \setminus x_1) \cup (f_2 \setminus x_2)) = e_1 \cap (f_1 \cup f_2 \setminus x_1 x_2) \). Hence, \( e_1 \cap (f_1 \cup f) \subseteq (e_1 \cap (f_1 \cup f_2)) \setminus x_j \) where \( \{i, j\} = \{1, 2\} \). Since \( e_1 \cap (f_1 \cap f_2) \) is monomial in \( B \), \( e_1 \cap f \subseteq e_1 \cap f_2 \) or \( e_1 \cap f_2 \subseteq e_1 \cap f \). But \( x_2 \notin (e_1 \cap f_2) \setminus (e_1 \cap f) \), then \( e_1 \cap f \not\subseteq e_1 \cap f_2 \). Now, if \( (f_1 \cap f) \subseteq e_1 \cap f_2 \), then \( f \subseteq (f_1 \cup f_2) \cap (f_1 \cap f) \cup (f_2 \cap f) \subseteq (e_1 \cup f_2) \cup (f_2 \cup f) \subseteq e_1 \cup f_2 \). So, \( f \subseteq (e_1 \cap f) \cup f_2 \subseteq (e_1 \cap f_2) \cup f_2 \subseteq f_2 \). But \( x_2 \notin f_2 \setminus f \), a contradiction. Hence, there is \( y_1 \in (f_1 \cap f) \setminus (e_1 \cup f_2) \). Similarly there is \( y_2 \in (f_2 \cap f) \setminus (e_1 \cap f_1) \). Consequently, the matrix

\[
\begin{pmatrix}
x_1 & x_2 & y_1 & y_2 \\
f_1 & 1 & 0 & 1 \\
f_2 & 0 & 1 & 0 \\
e_1 & 1 & 1 & 0 \\
f & 0 & 0 & 1 
\end{pmatrix}
\]

is a 4-cycle. A contradiction, therefore \( \mathcal{C} \) has the strong persistence property.

Example 3 ([1]). Let \( \mathcal{C}_0 \) be the clutter with vertex set \( \{x_1, \ldots, x_6\} \) whose edges are \( x_1 x_2 x_3, x_1 x_2 x_4, x_1 x_3 x_5, x_1 x_4 x_6, x_1 x_5 x_6, x_2 x_3 x_6, x_2 x_4 x_5, x_2 x_5 x_6, x_3 x_4 x_5 \) and \( x_3 x_4 x_6 \). \( \mathcal{C}_0 \) is an unmixed shellable clutter. But \( (I(\mathcal{C}_0))^2 : I(\mathcal{C}_0) \neq I(\mathcal{C}_0)^2 \), then \( \mathcal{C}_0 \) has no the strong persistence property.

Definition 9. The cone over a clutter \( \mathcal{C} \), denoted by \( \mathcal{C} x \), is the clutter whose vertex set is \( V(\mathcal{C}) \cup \{x\} \) and edge set \( \{f \cup \{x\} \mid f \in E(\mathcal{C})\} \), where \( x \) is a new vertex.

Proposition 5. \( \mathcal{C} \) has the strong persistence property if and only if \( \mathcal{C} x \) has the strong persistence property.
Proof. \(\Rightarrow\) If \(m = x^om' \in (I(C)x)^{k+1}: I(C)x)\) with \(gcd(m', x) = 1\), then \(\tilde{f}m \in I(C)x)^{k+1}\) for \(f \in E(C)\). Furthermore, \(\tilde{f} = \tilde{g}x\) with \(g \in E(G)\) then \(x^{k+1} | \tilde{g}xm\) implying \(x^k | m\). Thus, \(\alpha \geq k\) and \(\tilde{g}m' \in I(C)x)^{k+1}\). Hence \(m' \in I(C)x)^{k+1}\). \(i, e., m' = \ell f_1\cdots f_k\) where \(f_i \in E(C)\). Therefore, \(m = x^o\ell f_1\cdots f_k = x^o-k\ell (f_1x)\cdots(f_kx), \) so \(m \in I(C)x)^{k}\).

\(\Leftarrow\) If \(m \in (I(C)x)^{k+1}: I(C)x)\), then \(f m = g_1\cdots g_{k+1}\) for \(f \in I(C)\) and \(g_i \in I(C)\). Thus, \((fx)(mx^k) = \ell(xg_1)\cdots(xg_{k+1}) \in I(C)x)^{k+1}\). So, \(mx^k \in (I(C)x)^{k+1}: I(C)x)\) = \(I(C)x)^{k}\). Hence, \(mx^k = \ell(f_1x)\cdots(f_kx)\) for \(f_i \in I(C)\). Therefore \(m \in I(C)x)^{k}\).

\(\square\)

**Proposition 6.** \(C\) has the persistence property if and only if \(Cx\) has the persistence property.

Proof. If \(Q_1,\ldots, Q_r\) is the monomial minimal primary decomposition of \(I(C)x\) and \(Q'_i = R[x] \cdot Q_i\), then \(Q'_1,\ldots, Q'_r, (x^k)\) is the monomial minimal primary decomposition of \(I(C)x)^{k}\). Hence, \(\text{Ass}(I(C)x)^{k}) = \text{Ass}(I(C)x) \cup \{(x)\}\).

\(\square\)

**Proposition 7.** \(C = (V, E)\) has the strong persistence property if and only if \(C' = (V, E')\) has the strong persistence property, where \(E' = \{f \setminus \cap_{g \in E} g | f \in E\}\).

Proof. Set \(A = \cap_{g \in E} g\). By induction on \(k = |A|\). If \(k = 0\), then \(C = C'\). Now if \(k \geq 1\) and \(x \in A\), then \(C = C_1x\) where \(C_1 = C \setminus x\). So, by induction hypothesis \(C_1\) has the strong persistence property if and only if \(C'\) has the strong persistence property. Therefore, we obtain the result by Proposition 5.

\(\square\)

**Proposition 8.** A clutter \(C\) with 3 edges has the strong persistence property.

Proof. We assume \(E(C) = \{f_1, f_2, f_3\}\) and \(V(X) = \{x_1, \ldots, x_n\}\). By Proposition 7, we can suppose that \(f_1 \cap f_2 \cap f_3 = \emptyset\). If \(C\) is not connected, then it has a component with one edge. Hence, by Corollary 4 and Theorem 7, \(C\) has the strong persistence property. Now, we assume that \(C\) is connected. If \(f_i \cap f_j = \emptyset\) for some \(i \neq j\), then \(C\) has the strong persistence property by Lemma 2. Consequently, we suppose \(a_{ij} = f_i \cap f_j \neq \emptyset\) for \(i \neq j\). We set \(b_i\) such that \(f_i = a_{ij} \cup b_i \cup a_{ir}\) for \(i, j, r \in \{1, 2, 3\}\). So, each pair of \(b_1, b_2, b_3, a_{12}, a_{13}, a_{23}\) are disjoint.

We take \(m \in (I^{k+1}: I)\) where \(I = I(C)\), then \(mf_1 = \ell f_1 f_2 f_3\) with \(a_1 + a_2 + a_3 = k + 1\). If \(a_1 > 0\), then \(m \in I^k\). Now, if \(a_1 = 0\), then \(b_1 \mid \ell\) since \(b_1, f_1, f_3\) are disjoint pairs. This implies, \(\ell = b_1\ell'\) and \(ma_{12}a_{13} = \ell' f_2 f_3\). If \(a_2 = 0\), then \(\ell' = u_1a_{12}\) and \(m = u_1 v_2 f_3\) where \(f_3 = a_2 a_{13}\). Thus, \(m \in I^k\). Similarly if \(a_3 = 0\), then we suppose \(a_1 \neq 0\) and \(a_2 \neq 0\). Consequently \(m \ell' (b_2 b_3 a_{23}) f_2 f_3 \ell' = f_2 f_3\), implying \(a_{23} \mid m\) and \(b_2 b_3 f_3 \mid m\). Similarly, we can assume \(a_{23} \mid m\) and \(b_2 b_3 \mid m\). Hence, \((a_1 a_{12} a_{13}) f_2 f_3 \ell' = f_2 f_3\), implying \(a_{23} \mid m\) and \(b_2 b_3 f_3 \mid m\). Therefore, \(m \in I^{k+1} \subseteq I^k\).

\(\square\)

**Proposition 9.** If \(X\) is a set \(A \subseteq X\) and \(x \notin X\), then the clutter \(C\) whose edge set is \(\{X\} \cup \{xx_i | x_i \in A\}\) has the strong persistence property.
Proof. We set \( A = \{ x_1, \ldots, x_r \} \), \( f_0 = X \) and \( f_1 = \{ x, x_3 \} \). Since \( C \) is clutter, \( r > 1 \). We take \( m \in (I^{k+1}; I) \) where \( I = I(C) \), then \( m f_i = \ell_i f_0 + \ell_1 f_2 f_3 \ldots f_n \), where \( \sum_{i=0}^{r} \alpha_{ji} = k + 1 \) if \( \alpha_{0i} = 0 \) for each \( i \neq 1 \), then \( m \in (J^{k+1}; J) \), where \( J = (f_1, \ldots, f_r) \). But \( J \) is an edge ideal of a graph. So, by Theorem 2, \( m \in I^k \). Thus, we can assume \( \alpha_{01} > 0 \) and we take \( \alpha_1 = \alpha_{11} \). If \( \alpha_1 = 0 \) and \( x \notin E \), then \( x^{k-\alpha_0} \parallel m \) and \( \overline{f_0}^{\alpha_0} \parallel m \), since

\[
m = \ell_1 F_0^{\alpha_0} F_2^{\alpha_1} \ldots F_n^{\alpha_n} \quad \text{and} \quad \overline{f_0} = f_2^{\alpha_2} \ldots f_n^{\alpha_n}.
\]

So, \( x^{k-\alpha_0+1} \parallel m f_j \) and \( \overline{f_0}^{\alpha_0} \parallel m f_j \) for \( j \neq 1 \). Hence, \( m f_j \notin I^{k+1} \) a contradiction. Now if

\[
\alpha_1 \neq 0 \quad \text{or} \quad x \in \ell_1, \quad \text{then} \quad m = \ell_1 f_0 f_2 f_3 \ldots f_n \quad \text{or} \quad m = a b_1^{\alpha_0-1} f_1 f_2 \ldots f_n, \quad \text{where} \quad \ell_1 = xa \quad \text{and} \quad f_0 = x_1 b. \quad \text{Therefore} m \in I^k.
\]

Theorem 9. If \( I \) is a squarefree monomial ideal in \( K[x_1, x_2, x_3, x_4] \), then \( I \) has the strong persistence property.

Proof. Let \( C \) be the clutter associated to \( I \). By Proposition 8 and Theorem 7 we can assume that \( |E(C)| > 3 \) and \( C \) has no edges of cardinality 1. If \( C \) has only edges of cardinality 3, then \( 4 \leq |E(C)| \leq \binom{\alpha}{2} = 4 \). Hence, \( C \) is a complete clutter, implies \( C \) is a base set of a polymatroid. Consequently, by [1, Proposition 2.4] \( C \) has the strong persistence property. If \( C \) has only one edge of cardinality 2, then \( |E(C)| \leq 3 \). A contradiction, so there are \( f_1, f_2 \in E(C) \) such that \( |f_1| = |f_2| = 2 \). By Theorem 2 we can suppose \( f = \{ x_1, x_2, x_3 \} \in E(C) \). So, if \( f' \in E(C) \setminus \{ f \} \), then \( x_4 \in f' \). Hence, we can assume \( f_1 = \{ x_1, x_4 \} \) and \( f_2 = \{ x_2, x_3 \} \). Thus, if \( f' \in E(C) \setminus \{ f_1, f_2 \} \), then \( f' = \{ x_3, x_4 \} \). Therefore, by Proposition 9, \( C \) has the strong persistence property.

Corollary 7. If \( I \subseteq K[x_1, \ldots, x_n] \) is a squarefree monomial ideal without the strong persistence property, then \( n \geq 5 \) and there is \( k \geq 3 \) such that \( (I^k; I) \neq I^{k-1} \).


Definition 10. Let \( C = (V, E) \) be a clutter with \( x \in V \), the deleting of \( x \) is the clutter \( C \setminus x \) with vertex set \( V \setminus \{ x \} \) and edge set \( \{ f \in E \mid x \notin f \} \). Furthermore, the contraction of \( x \) is the clutter \( C / x \) with vertex set \( V \setminus \{ x \} \) and whose edges are \( f \setminus \{ x \} \) with \( f \in E \) and there is not \( f' \in E \) such that \( f' \setminus \{ x \} \subset f \setminus \{ x \} \).

Example 4. We consider the clutter \( C \) with vertex set \( V(C_0) \cup \{ x \} \) and edge set \( E(C_0) \cup \{ xx_1 \} \), where \( C_0 \) is the clutter in Example 3. By Theorem 7, \( I(C) \) has the strong property but \( C \setminus x = C_0 \) has not the strong persistence property.

Proposition 10. Let \( C \) be a clutter and \( x \in V(C) \). If \( C \) has the (strong) persistence property, then \( C / x \) has the (strong) persistence property.
Proof. We set \( E(C) = \{ f_1, \ldots, f_r \} \). We can suppose \( \{ f_i \mid x \in f_i \} = \{ f_1, \ldots, f_r \} \) and \( \{ f_i \mid x \in f_i \} = \{ f_{r+1}, \ldots, f_J \} \). We define \( f'_i = f_i \setminus \{ x \} \) for \( i \leq r_2 \) and \( A = \bigcup_{i \leq r_2} f'_i \). Also, we set \( I = I(C \setminus x) \) and \( J = I(C) \). Thus, \( f'_1, \ldots, f'_r \) are the edges of \( C \setminus x \) and \( f'_1 \) for \( r_1 + 1 \leq i \). Furthermore, if \( i > r_2 \), then \( f'_i \subseteq f_i \) for some \( i \). So, for each \( 1 \leq i \leq r \) there is \( j \leq r_2 \) such that \( f'_j \mid f_i \). Consequently, if \( m \in G(J^k) \), then there is \( m' \in G(I^k) \) such that \( m' | m \). We take \( \mathcal{L} = (x_1, \ldots, x_r) \) where \( z_j = x_j^{\beta_j} \). Hence if \( \mathcal{L} \) is an \( I^k \)-minimal ideal, then \( J^k \subseteq \mathcal{L} \). Furthermore, \( \mathcal{G}(r(\mathcal{L})) = \{ x_1, \ldots, x_r \} \subseteq A \) since \( \mathcal{L} \) is \( I^k \)-minimal. Now, we suppose \( \mathcal{L} \) is \( J^k \)-minimal and \( \mathcal{G}(r(\mathcal{L})) \subseteq A \). If \( m \in G(I^k) \), then \( m = \prod_{j=1}^{a_{x_1}} \prod_{j=2}^{a_{x_2}} \prod_{j=r}^{a_{x_r}} \) with \( a_1 + \ldots + a_{r_2} = k \). So, \( x''m = \prod_{j=1}^{a_{x_1}} \prod_{j=2}^{a_{x_2}} \prod_{j=r}^{a_{x_r}} \) is \( I^k \), where \( \alpha = a_1 + \ldots + a_{r_1} \). Thus, \( z_j \mid x''m \) for some \( 1 < s \). Since \( x \notin A \), \( \gcd(x, z_j) = 1 \), and \( z_j \mid m \). Therefore \( I^k \subseteq \mathcal{L} \). Now, we will prove that \( \mathcal{L} \) is an \( I^k \)-minimal ideal if and only if \( \mathcal{L} \) is an \( J^k \)-minimal ideal and \( \mathcal{G}(r(\mathcal{L})) \subseteq A \). Assume \( \mathcal{L} \) is \( I^k \)-minimal so, \( r(\mathcal{L}) \subseteq A \). If \( \mathcal{L} \) is not \( J^k \)-minimal, then there is \( \mathcal{L}' \) such that \( J^k \subseteq \mathcal{L}' \subseteq \mathcal{L} \) and \( r(\mathcal{L}') \subseteq r(\mathcal{L}) \subseteq A \). Consequently, \( I^k \subseteq \mathcal{L}' \). A contradiction, therefore \( \mathcal{L} \) is \( J^k \)-minimal. Now suppose \( \mathcal{L} \) is \( J^k \)-minimal and \( \mathcal{L} \) is not \( I^k \)-minimal, then there is \( \mathcal{L}' \) such that \( I^k \subseteq \mathcal{L}' \subseteq \mathcal{L} \). This implies \( J^k \subseteq \mathcal{L}' \), a contradiction, since \( \mathcal{L} \) is \( J^k \)-minimal.

Hence, \( \text{Ass}(I^k) = \{ P \in \text{Ass}(J^k) \mid G(P) \subseteq A \} \) for each \( k \). Since \( \mathcal{L} \) has the persistence property, if \( P \in \text{Ass}(J^k) \), then \( P \in \text{Ass}(I^{k+1}) \) and \( G(P) \subseteq A \). Thus, \( P \in \text{Ass}(I^{k+1}) \). Therefore, \( \mathcal{L} \) has the persistence property.

(Strong). Now, we set \( m \in (I^{k+1}I) \). If \( 1 \leq i \leq r_2 \), then \( m = \ell_i f_1^{\alpha_1} \ldots f_r^{\alpha_r} \) where \( \ell_i \in \text{Mon}(R) \) and \( \alpha_1 + \ldots + \alpha_{r_2} = k + 1 \). We take \( u_i = \alpha_1 + \ldots + \alpha_{r_1} \). If \( i \leq r_1 \), then
\[
\alpha_i + \ldots + \alpha_{r_2} = k + 1.
\]

Now if \( r_1 + 1 \leq i \leq r_2 \), then \( x^{k+1}m = x^{k+1}m = x^{k+1}f_1^{\alpha_1} \ldots f_{r_2}^{\alpha_{r_2}} \). Finally if \( r_2 + 1 \leq i \leq r \), then there exist \( j \leq r_1 \) such that \( f'_j \mid f_i \). So,
\[
\alpha_i + \ldots + \alpha_{r_2} = k + 1.
\]

Consequently, \( x^{k+1}m \in (J^{k+1}, J) = J^k \). This implies \( x^{k+1}m = \ell f_1^{\rho_1} \ldots f_r^{\rho_r} \) with \( \ell \in \text{Mon}(R) \) and \( \beta_1 + \ldots + \beta_r = k \). Since \( x \notin I \), \( x \mid \ell f_j \) for \( j \geq r_1 + 1 \), \( x \mid \ell \), where \( w = k + 1 - (\beta_1 + \ldots + \beta_r) \). Therefore, \( \ell = x^w \ell' \) where \( \ell' \in \text{Mon}(R) \) and \( m = \ell f_1^{\beta_1} \ldots f_r^{\beta_r} \).

\[ \square \]

Remark 2. The converse affirmation of Proposition 10 is not true. We take \( C_0 \) as in Example 3. So, \( C_0 \neq \{ x_i \} \) is a simple graph for each \( i \). Hence, by Theorem 3, \( C_0 \neq \{ x_i \} \) has the strong persistence property.

Definition 11. Let \( C = (V, E) \) be a clutter and \( \sigma \in S_V \) a permutation. We consider the clutter \( \sigma(C) = (V, E') \) where \( E' = \{ x_{\sigma(i_1)} \ldots x_{\sigma(i_r)} \mid x_{i_1}, \ldots, x_{i_r} \in E \} \).

Proposition 11. If \( C \) has the strong persistence property and \( \sigma \in S_V(C) \), then \( \sigma(C) \) also has the strong persistence property.

Proof. We take a morphism of \( k \)-algebras \( \phi : R = K[x_1, \ldots, x_n] \rightarrow R \) given by \( \phi(x_i) = x_{\sigma(i)} \).

Hence, \( \phi \) is an automorphism of \( R \), with \( \phi(I(C)) = I(\sigma(C)) \). Therefore, \( I(C) \) and \( I(\sigma(C)) \) are isomorphic.
5 The symbolic strong persistence property

In this section we study some properties of the strong persistence property in a general ring. Furthermore, we introduce the symbolic strong persistence property and we prove that the strong persistence property implies the symbolic strong persistence property.

**Theorem 10.** An ideal $I$ has the strong persistence property if and only if $(I^t:I^s) = I^{t-s}$ for all $s \leq t$.

**Proof.** We proceed by induction on $s$. For $s = 1$ we recover the strong persistence property. Now, we take $a \in (I^t:I^{s+1})$ with $t \geq s+1$ and $x \in I$, then $axb \in I^t$ for all $b \in I^s$. Hence $ax \in (I^t:I^s)$. By induction hypothesis $ax \in I^{t-s}$. Consequently $a \in (I^{t-s}:I)$ and, by induction, $a \in I^{t-s-1}$. Therefore $(I^t:I^{s+1}) = I^{t-s-1}$.

**Corollary 8.** If $I$ has the strong persistence property, then $I^t$ has the strong persistence property.

**Proof.** By Theorem 10 $(I^{kt}:I^t) = I^{kt-t} = I^{(k-1)t}$ for all $k \geq 1$. Therefore, $I^t$ has the strong persistence property.

By [4] normal ideals in an integer domain satisfy $(I^t:I^s) = I^{t-s}$ for all $s \leq r$. Hence, by Theorem 10 a normal ideal has the strong persistence property, but the converse affirmation is not true.

**Example 5 ([5]).** Let $G$ be a simple connected graph, the $I(G)$ has the strong persistence property but if $V(G) = \{x_1, x_2, x_3, x_4, x_5, x_6, x_7\}$ and $E(G) = \{x_1x_2, x_2x_3, x_1x_3, x_3x_4, x_4x_5, x_5x_6, x_6x_7, x_5x_7\}$, then $I(G)$ is not normal.

**Definition 12.** Let $P_1, \ldots, P_r$ be the minimal primes of $I$. The $i$-th symbolic power of $I$ is $I^{(i)} = q_1 \cap \cdots \cap q_i$, where $q_i$ is the $P_i$-primary component of $I^n$.

**Remark 3.** $I^{(i)} \subseteq (I^{(i+1)}:I^{(1)})$ for each $i$.

**Definition 13.** $I$ has the symbolic strong persistence property if $(I^{(i+1)}:I^{(1)}) = I^{(i)}$ for each $i$.

**Theorem 11.** Strong persistence property implies the symbolic strong persistence property.

**Proof.** Let $\text{Min}(I) = \{P_1, \ldots, P_r\}$ be the set of minimal primes containing $I$. We take $I^d = Q_{1d} \cap \cdots \cap Q_{sd}$ a minimal primary decomposition of $I^d$ for each $d$. We can suppose that there exists $r_d \leq s_d$ such that $r(Q_{id}) \in \text{Min}(I)$ if and only if $i \leq r_d$. Now for $j > r_{k+1}$, then $r(Q_{j,k+1})$ is not minimal. Consequently, $r(Q_{j,k+1}) \notin r(Q_{i,k+1})$ with $i \leq r_{k+1}$. This implies $r(Q_{j,k+1}) \notin B$, where $B = \bigcup_{i \geq 1} r(Q_{i,k+1})$. Thus, there is $a_j \in r(Q_{j,k+1}) \setminus B$. So, $b_j = a_j^{s_j} \in Q_{j,k+1}$ for some $s_j$. Hence, $b_j \in Q_{j,k+1} \setminus B$. Now, we take $a \in (I^{(k+1)}:I^{(1)})$, then $ax \in I^{(k+1)}$ for all $x \in I^{(1)}$. Consequently, if $c = \prod_{j \geq r_{k+1}} b_j$, then $axc \in I^{(k+1)}$ for all $x \in I$ since $I \subseteq I^{(1)}$. So, $ac \in I^k$, since $I$ has the strong persistence property. Furthermore, if $j > r_{k+1}$, then $b_j \notin r(Q_{ik})$ for $i \leq r_k$. Thus, $a \in Q_{ik}$ for $1 \leq i \leq r_k$, since $ac \in Q_{ik}$ and $Q_{ik}$ is primary. Therefore, $a \in I^k$.

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Proposition 12. An ideal $I$ has the symbolic strong persistence property if and only if $(I^{(r)}; I^{(s)}) = I^{(r-s)}$ for all $s \leq r$.

Proof. Similar to proof of Theorem 10. □

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References


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